# Lecture 22 I/O Performance and Checkpoints

EN 600.320/420/620

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#### The I/O Crisis in HPC

In a world where FLOPS is the commodity......

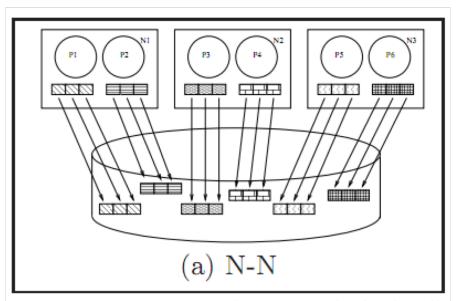
....Disk I/O often limits performance

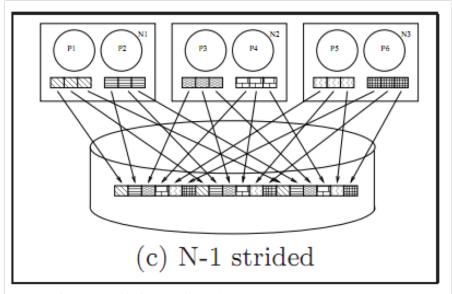
- Any persistent data must make it off the supercomputer
  - To magnetic or solid state storage
- Storage is not as connected to the high-speed network as compute
  - Because it needs to be shared with other computers
  - Because it doesn't add to TOP500 benchmarks



#### Where does the I/O Come From?

- Checkpointing!
  - And, writing output from simulation (which is checkpointing)
- Checkpoint workload
  - Every node node writes local state to a shared file system
  - Using POSIX calls (FS parallelized) or MPI I/O









# Why Checkpointing

- At scale failures occur inevitably
  - MPI synchronous model means that a failure breaks the code
  - Lose all work since start (or restart)
- Each checkpoint provides a restart point
  - Limits exposure, loss of work to last checkpoint
- By policy, all codes that run at scale on supercomputers MUST checkpoint!
  - HPC centers want codes to do useful work



## **Checkpoint Approaches**

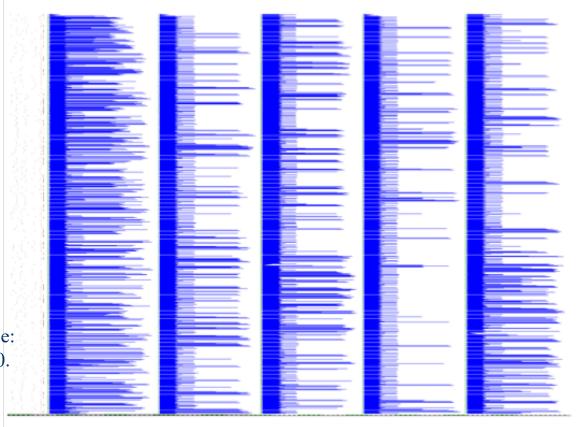
- Automatic: store contents of memory and program counters
  - Brute force, large data, inefficient
  - But easy, no development effort
  - New interest in this approach with the emergence of VMs and containers in HPC.
- Application specific: keep data structures and metadata representing current progress. Hand coded by developer.
  - Smaller, faster, preferred, but tedious.
  - Almost all "good" codes have application specific checkpoints



# **A Checkpoint Workload**

- IOR benchmark
  - Each node transfers 512 MB
- Barriers
  - How much parallelism?
- What effects?

M. Uselton et al. Parallel I/O Performance: From Events to Ensembles. IPDPS, 2010.

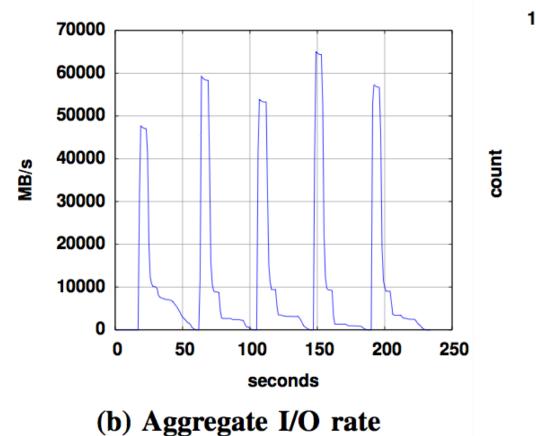


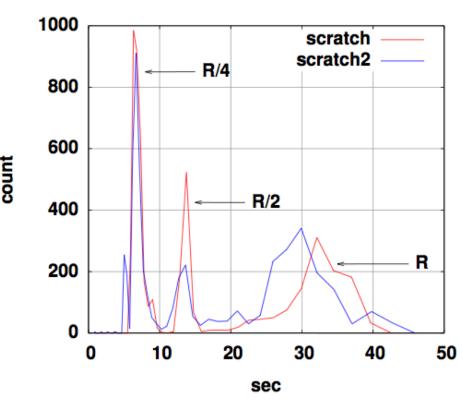
(a) I/O trace diagram



## I/O Rates and PDF

What features do you observe?





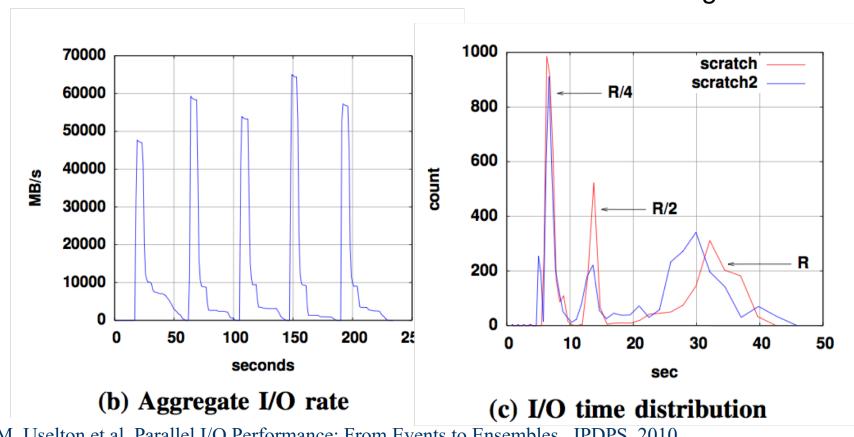
(c) I/O time distribution



M. Uselton et al. Parallel I/O Performance: From Events to Ensembles. IPDPS, 2010.

## I/O Rates and PDF

- What features do you observe?
  - Lagging processes = not realizing peak I/O performance
  - Harmonics in I/O distribution = unfair resource sharing





M. Uselton et al. Parallel I/O Performance: From Events to Ensembles. IPDPS, 2010.

## **Statistical Observations**

#### Order statistics

Fancy way of saying, the longest operation dominates overall performance

#### Law of large numbers

- I don't think that they make this analysis cogent
- It's right, but Gaussian distribution is not what matters
- A better, intuitive conclusion is

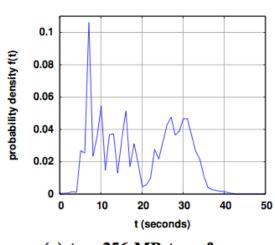
#### (RB interprets) smaller files are better

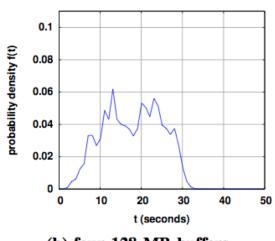
- The worst case slow down on a smaller transfer takes less absolute time than on a large transfer
- As long as transfers are "big enough" to amortize startups costs

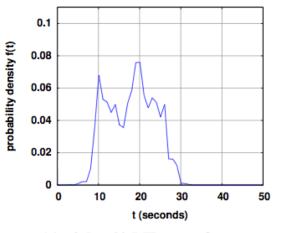


# **Smaller Files Improve Performance**

- Non-intuitive
  - Smaller operations seems like more overhead
  - But, a property of statistical analysis
- Smaller better as long as fixed costs are amortized
  - Obviously, 1 byte is too small







(a) two 256 MB transfers

(b) four 128 MB buffers

(c) eight 64 MB transfers

Figure 2: IOR 512 MB transfer using 1024 processors where: a) 512 MB written via two 256MB write() calls. b) Four calls (128MB). c) Eight calls (64MB). Note that the distributions become progressively narrower and more Gaussian.



# **The Checkpoint Crisis**

As HPC codes get larger, I/O becomes more critical

- Some observations
  - Checkpoint to protect against failure
  - More components increase failure probability
  - FLOPs grows faster than bandwidth
- Conclusion
  - Must take slower checkpoints more often
  - Eventually you will get no constructive work done between checkpoints
- Mitigation (just delaying the problem)
  - Burst buffers: fast (SSD) storage in high-speed network
  - Observe the checkpoint persistence is shorter than needed for output/analysis data

